SI2-SSI: Collaborative Research: A Sustainable Infrastructure for Performance, Security, and Correctness Tools

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Awards: 1450273, 1449918

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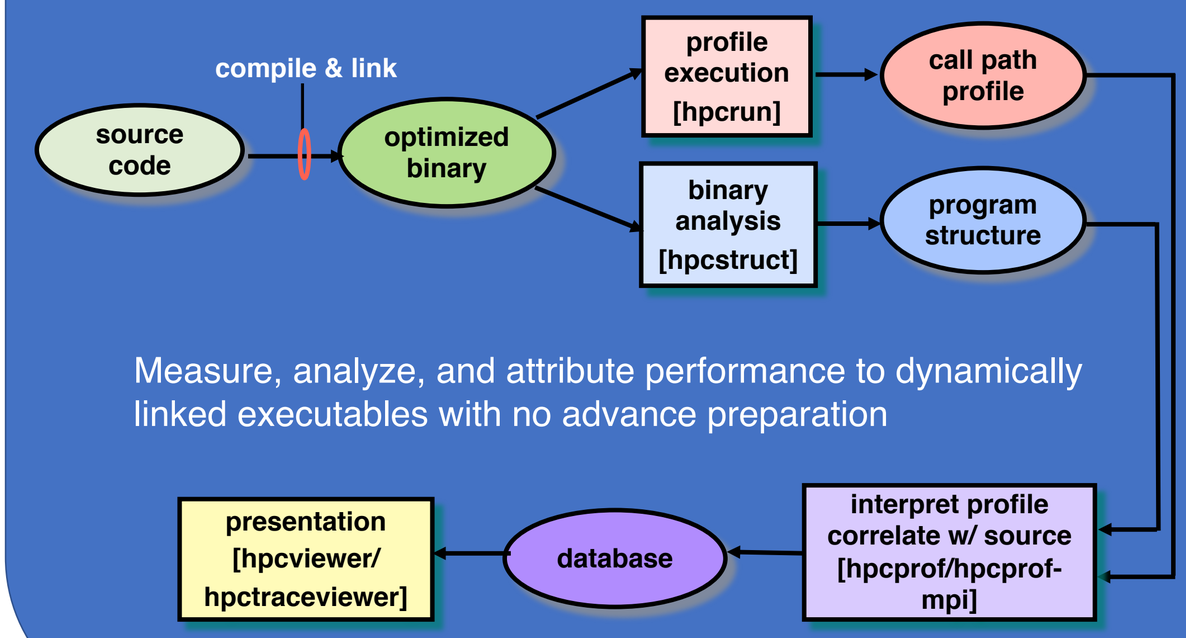
HPCToolkit: Performance Measurement, Attribution, and Analysis Tool Suite

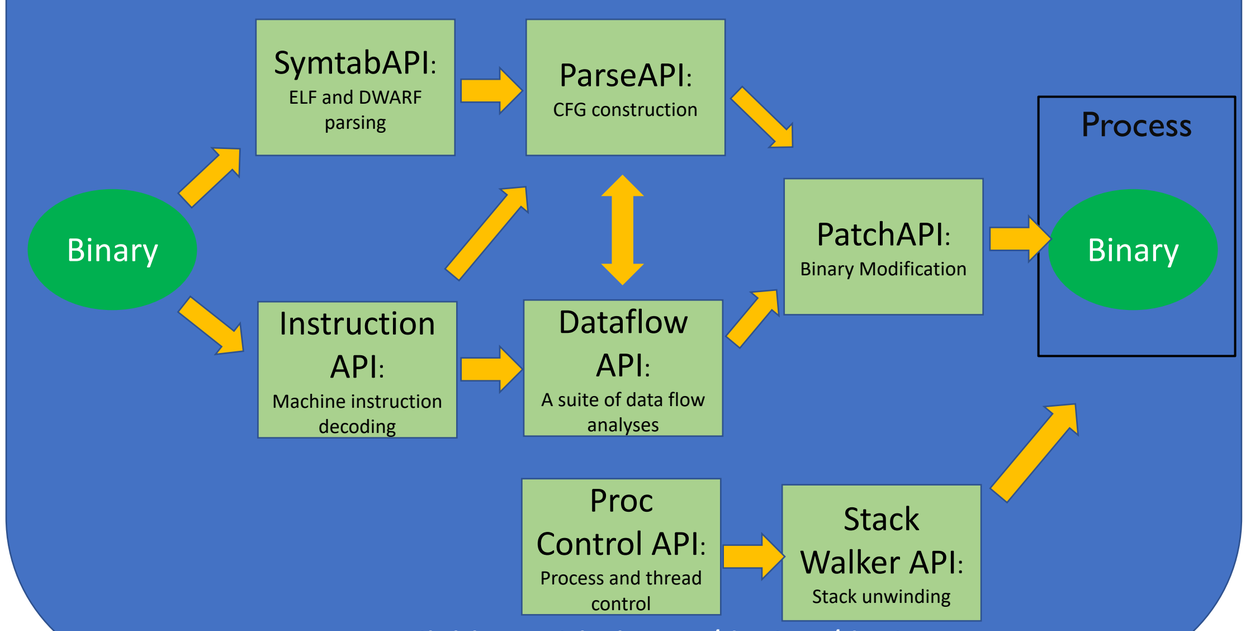


Measure, analyze, and attribute performance to dynamically linked executables with no advance preparation

Available at hpctoolkit.org

Dyninst: Binary Analysis and Instrumentation Tool Suite

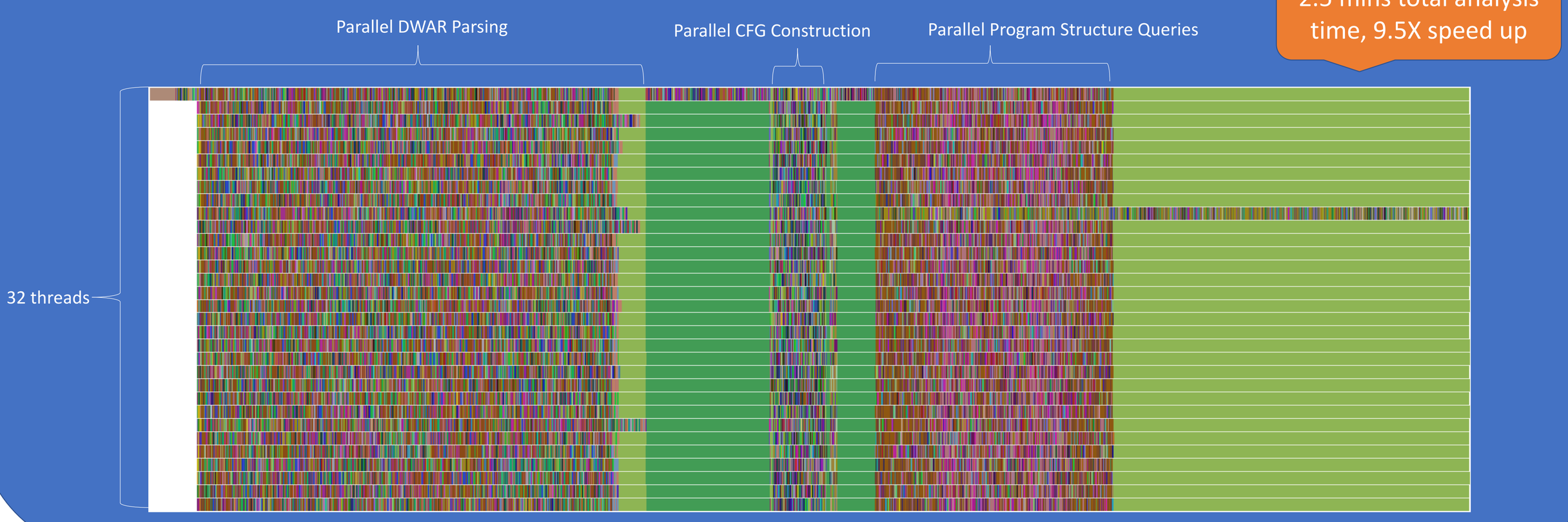
Component libraries for binary structure and code analysis, dynamic instrumentation, and static binary rewriting:



Available at github.com/dyninst/dyninst

Analyze Large Scale Binaries

- Added multi-threading to Dyninst's ParseAPI and SymtabAPI [1]
- HPCToolkit's hpcstruct uses Dyninst's ParseAPI to analyze loop nesting, source line mapping, and function inlining
- A trace view of hpcstruct analyzing a 8.2GB shared library from TensorFlow



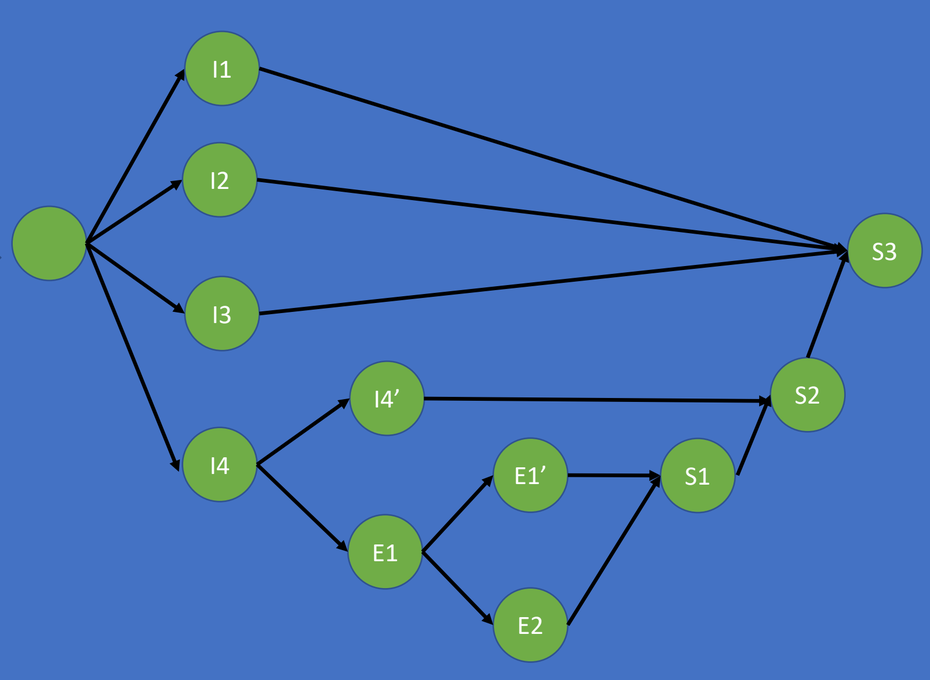
[1] Parallelizing Binary Code Analysis, Xiaozhu Meng, Jonathon M. Anderson, John Mellor-Crummey, Mark W. Krentel, Barton P. Miller, Srđan Milaković, https://arxiv.org/abs/2001.10621

Data Race Detection for OpenMP Programs

- Instrument memory reads and writes using Dyninst to track happens-before relation
- Reduce state information maintained by leveraging OpenMP semantics

```
#pragma omp parallel for
for (int i = 0; i < 4; ++i) {  // I1 – I4
    if (i < 3) {
        int local = shared;
        // Processing local
    } else {
        #pragma omp task // E1
        {
            int local = shared;
            #pragma omp task // E2
            {
                shared *= 2;
            }
            #pragma omp taskwait
        }
    }
}
```

Covert to OpenMP task graph



Use static binary analysis to reduce instrumentation:

- Avoid instrumentation for accesses to the same variable by the same task in the same synchronization interval
- Avoid instrumentation for accesses to read-only memory locations

Basic algorithm execution trace
1. I1, I2, and I3 perform concurrent read
-> Record access history: I1, I2, and I3 read `shared`
2. E1 and E2 perform concurrent read
-> Record access history: E1 and E2 read `shared`
3. E2 performs a concurrent write
-> Data race
Five task IDs are recorded for `shared`

New algorithm execution trace
1. I1, I2, and I3 perform concurrent read
-> Record access history: I1 and I2 read `shared`; no need to record I3
2. E1 and E2 perform concurrent read
-> no need to record E1 or E2
3. E2 performs concurrent write
-> Data race
Two task IDs are recorded for `shared`